

ON UNIFORM f -VECTORS OF CUTSETS IN THE TRUNCATED BOOLEAN LATTICE

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Let $[n] = \{1, 2, \dots, n\}$ and let $2^{[n]}$ be the collection of all subsets of $[n]$ ordered by inclusion. $\mathcal{C} \subseteq 2^{[n]}$ is a *cutset* if it meets every maximal chain in $2^{[n]}$, and the *width* of $\mathcal{C} \subseteq 2^{[n]}$ is the minimum number of chains in a chain decomposition of \mathcal{C} . Fix $0 \leq m \leq l \leq n$. What is the smallest value of k such that there exists a cutset that consists only of subsets of sizes between m and l , and such that it contains exactly k subsets of size i for each $m \leq i \leq l$? The answer, which we denote by $g_n(m, l)$, gives a lower estimate for the width of a cutset between levels m and l in $2^{[n]}$. After using the Kruskal–Katona Theorem to give a general characterization of cutsets in terms of the number and sizes of their elements, we find lower and upper bounds (as well as some exact values) for $g_n(m, l)$.

1. Introduction

Let $2^{[n]}$ be the *Boolean lattice* of order n , that is the lattice of all subsets (often called *nodes*) of $[n] = \{1, 2, \dots, n\}$ ordered by inclusion. For $0 \leq m \leq n$ we define the m -th *level set* $\binom{[n]}{m}$ of $2^{[n]}$ as the set of all subsets of size m . The f -vector (or *profile*) $\mathbf{f} = (f_0, f_1, \dots, f_n)$ of a collection of subsets $\mathcal{A} \subseteq 2^{[n]}$ is defined by $f_m = |\mathcal{A}_m|$ where $\mathcal{A}_m = \mathcal{A} \cap \binom{[n]}{m}$ and $0 \leq m \leq n$.

A collection of $l+1$ subsets $A_0 \subset A_1 \subset \dots \subset A_l$ in $2^{[n]}$ is called a *chain* of length l . A *maximal chain* in $2^{[n]}$ is one that has length n . A collection of w nodes with the property that none of them contains another is called an *antichain* of size w . The *length* and the *width* of a collection of subsets $\mathcal{A} \subseteq 2^{[n]}$ are defined as the length of the longest chain and the size of the largest antichain in \mathcal{A} , respectively.

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A *cutset* in $2^{[n]}$ is defined as a collection of subsets $\mathcal{C} \subseteq 2^{[n]}$ which intersects all maximal chains. Trivially, every collection \mathcal{C} which contains \emptyset or $[n]$ is a cutset. In [3] we proved that for $n \geq 2$, the width of a cutset in the Boolean lattice of order n which does not contain \emptyset or $[n]$ is greater than or equal to $n-1$, and that for $n \geq 3$ there exist cutsets of width $n-1$ in $2^{[n]}$. Thus, it is possible to construct a cutset in $2^{[n]}$ with f -vector $(0, \underbrace{n-1, n-1, \dots, n-1}_{n-1}, 0)$.

We then may ask for the smallest value of k for which there is a cutset in $2^{[n]}$ with f -vector $(0, \underbrace{k, k, \dots, k}_{n-1}, 0)$. The original goal of our work was to show that this value is $n-2$ (see [Corollary 3](#) below).

More generally, for $0 \leq m \leq l \leq n$ we define $g_n(m, l)$ to be the smallest value of k for which the $n+1$ -tuple (f_0, f_1, \dots, f_n) , defined by $f_i = k$ if $m \leq i \leq l$ and 0 otherwise, can be the f -vector of a cutset in $2^{[n]}$. Thus our goal above is then to find $g_n(1, n-1)$. Note that by symmetry we have $g_n(m, l) = g_n(n-l, n-m)$, so we may assume without loss of generality that $m \leq l \leq n-m$.

Before studying $g_n(m, l)$, we give a general characterization of f -vectors of cutsets in $2^{[n]}$. For a given profile $\mathbf{f} = (f_0, f_1, \dots, f_n)$ and integer m_0 with $0 \leq m_0 \leq n$, we construct a *canonical* collection of subsets $\mathcal{C}(\mathbf{f}, m_0)$, with the property that there is a cutset in $2^{[n]}$ with profile \mathbf{f} if and only if $\mathcal{C}(\mathbf{f}, m_0)$ is a cutset for some (or every) $0 \leq m_0 \leq n$. We then translate this qualitative criterion to a quantitative one: For a given $\mathbf{f} = (f_0, f_1, \dots, f_n)$ and $0 \leq m_0 \leq n$, we describe an easily computable value $q(\mathbf{f}, m_0)$, so that \mathbf{f} will be the profile of a cutset in $2^{[n]}$ exactly when $f_{m_0} \geq q(\mathbf{f}, m_0)$ for some (or every) $0 \leq m_0 \leq n$. These characterizations, which we present in [Section 2](#), are essentially due to Daykin [7] (for a correction see [5] and then [4]), though we follow a treatment which is more suitable for our purposes.

We can then determine the values of $g_n(m, l)$ for $l \leq m+2$. Namely, we prove the following.

Theorem 1.1. *Let n be a positive integer.*

1. $g_n(m, m) = \binom{n}{m}$ for every integer $0 \leq m \leq n$.
2. $g_n(m, m+1) = \binom{n-1}{m}$ for every integer $0 \leq m \leq n-1$.
3. $g_n(m, m+2) = \sum_{j=0}^m \binom{n-2j-2}{m-j}$ for every integer $0 \leq m \leq n/2-1$.

Next, viewing m as fixed and $n \gg m$ (i.e., for all $n > n_0 = n_0(m)$), we develop upper and lower bounds for $g_n(m, l)$.

Theorem 1.2. Suppose that m and n are non-negative integers and $n \gg m$. Then

1. $\binom{n-2}{m} < g_n(m, l) \leq \sum_{j=0}^m \binom{n-2j-2}{m-j}$ for every integer $m+2 \leq l \leq n-m-1$,
and
2. $\binom{n-3}{m} < g_n(m, n-m) \leq \sum_{j=0}^m \binom{n-2j-3}{m-j}$.

For $m=1$ we then get the following results.

Corollary 1.3. Suppose that $n > 4$ and $1 \leq l \leq n-1$ are integers. Then

$$g_n(1, l) = \begin{cases} n & \text{if } l = 1 \\ n-1 & \text{if } 2 \leq l \leq n-2 \\ n-2 & \text{if } l = n-1 \end{cases}$$

For $2 \leq m \ll n$, [Theorems 1 and 2](#) give the “numerator” of the leading term of the m -binomial representation (see [section 2](#)) of $g_n(m, l)$. Namely, this value is equal to n if $l=m$, $n-1$ if $l=m+1$, $n-2$ if $m+2 \leq l \leq n-m-1$, and $n-3$ if $l=n-m$. It is striking that for a rather large range of values of l , $g_n(m, l)$ stays essentially unchanged.

We note that in [Theorem 2](#), the ratio of the upper bound to the lower bound is approximately $1 + \frac{m}{n}$, and thus the bounds are rather accurate as $n \gg m$.

Extremal problems regarding cutsets in the Boolean lattice have been the object of much study. For example see [\[8\]](#), [\[10\]](#), [\[12\]](#), [\[13\]](#), [\[15\]](#), [\[17\]](#), [\[18\]](#).

2. f -vectors of cutsets

Given a collection $\mathcal{B} \subseteq \binom{[n]}{m}$, the *shadow* and the *shade* of \mathcal{B} will be denoted by $\triangle \mathcal{B}$ and $\nabla \mathcal{B}$, respectively [\[1, Chapter 2\]](#), and are as usual defined by

$$\begin{aligned} \triangle \mathcal{B} &= \left\{ A \in \binom{[n]}{m-1} \mid A \subseteq B \text{ for some } B \in \mathcal{B} \right\}, \\ \nabla \mathcal{B} &= \left\{ A \in \binom{[n]}{m+1} \mid B \subseteq A \text{ for some } B \in \mathcal{B} \right\}. \end{aligned}$$

We order the elements of $\binom{[n]}{m}$ by the *squashed order* (also called the *colex* order) [\[1, Chapter 7\]](#), that is for $A, B \in \binom{[n]}{m}$, we say $A <_S B$ if the largest element of the symmetric difference of A and B is in B . For $1 \leq K \leq \binom{n}{m}$, we define the *initial collection* $\mathcal{F}_m(K)$ and the *last collection* $\mathcal{L}_m(K)$ at level m as the first and last K elements in the squashed order at level m , respectively.

In addition, if $K \leq 0$, then $\mathcal{F}_m(K) = \mathcal{L}_m(K) = \emptyset$. The squashed order has the property that the shadow of an initial collection at level m is an initial collection at level $m - 1$, and the shade of a last collection at level m is a last collection at level $m + 1$. The Kruskal–Katona Theorem ([14], [16] or [1, Chapter 7]) states that the size of the shadow of K nodes at level m is greater than or equal to the size of the shadow of $\mathcal{F}_m(K)$ and, equivalently, the size of their shade is greater than or equal to the size of the shade of $\mathcal{L}_m(K)$.

Let Ω_n denote the set of $n+1$ -tuples of integers (a_0, a_1, \dots, a_n) such that $0 \leq a_m \leq \binom{n}{m}$ for all $0 \leq m \leq n$. To see whether a given $\mathbf{f} \in \Omega_n$ is the profile of a cutset in $2^{[n]}$, we construct a collection of subsets $\mathcal{C} = \mathcal{C}(\mathbf{f}, m_0)$, called the *canonical* collection of subsets for profile \mathbf{f} and for level m_0 ($0 \leq m_0 \leq n$). As we show below, there is a cutset in $2^{[n]}$ with profile \mathbf{f} if and only if this canonical collection is a cutset for some (or every) m_0 .

Our construction is as follows. First we let $\mathcal{E}_0^\uparrow = \{\emptyset\}$, $\mathcal{C}_0^\uparrow = \mathcal{F}_0(f_0)$, and for $1 \leq m \leq n$ we recursively define $\mathcal{E}_m^\uparrow = \nabla(\mathcal{E}_{m-1}^\uparrow - \mathcal{C}_{m-1}^\uparrow)$ and $\mathcal{C}_m^\uparrow = \mathcal{L}_m(|\mathcal{E}_m^\uparrow|) - \mathcal{L}_m(|\mathcal{E}_m^\uparrow| - f_m)$. Then \mathcal{E}_m^\uparrow is a last collection at level m , and it is precisely the set of nodes from which there is a chain of length m to \emptyset which is disjoint from \mathcal{C}_i^\uparrow for all $0 \leq i \leq m-1$. Analogously, we let $\mathcal{E}_n^\downarrow = \{[n]\}$, $\mathcal{C}_n^\downarrow = \mathcal{L}_n(f_n)$, and for $0 \leq m \leq n-1$ we recursively define $\mathcal{E}_m^\downarrow = \Delta(\mathcal{E}_{m+1}^\downarrow - \mathcal{C}_{m+1}^\downarrow)$ and $\mathcal{C}_m^\downarrow = \mathcal{F}_m(|\mathcal{E}_m^\downarrow|) - \mathcal{F}_m(|\mathcal{E}_m^\downarrow| - f_m)$. This time \mathcal{E}_m^\downarrow is an initial collection at level m , and it is the set of nodes from which there is a chain of length $n-m$ to $[n]$ which is disjoint from \mathcal{C}_i^\uparrow for all $m+1 \leq i \leq n$. Finally, we define $\mathcal{C} = \mathcal{C}(\mathbf{f}, m_0) = (\cup_{i=0}^{m_0} \mathcal{C}_i^\uparrow) \cup (\cup_{i=m_0+1}^n \mathcal{C}_i^\downarrow)$.

We can easily see that $\mathcal{C} = \mathcal{C}(\mathbf{f}, m_0)$ is a cutset if and only if $\mathcal{E}_{m_0}^\uparrow \cap \mathcal{E}_{m_0}^\downarrow \subseteq \mathcal{C}_{m_0}$. Furthermore, the profile (c_0, c_1, \dots, c_n) of \mathcal{C} satisfies $c_m \leq f_m$ for every m , and if \mathcal{C} is not a cutset, then its profile is exactly \mathbf{f} .

For example, let $n = 5$. If $\mathbf{f} = (0, 2, 5, 6, 0, 0)$ then, for all $0 \leq m_0 \leq 5$, $\mathcal{C}(\mathbf{f}, m_0)$ becomes

$$\begin{aligned} &\{\{1\}, \{2\}, \{1, 3\}, \{2, 3\}, \{1, 4\}, \{2, 4\}, \{3, 4\}, \{1, 2, 5\}, \{1, 3, 5\}, \\ &\quad \{2, 3, 5\}, \{1, 4, 5\}, \{2, 4, 5\}, \{3, 4, 5\}\}. \end{aligned}$$

On the other, hand if $\mathbf{g} = (0, 2, 6, 5, 0, 0)$, then for $\mathcal{C}(\mathbf{g}, 5)$ we get

$$\begin{aligned} &\{\{1\}, \{2\}, \{1, 3\}, \{2, 3\}, \{1, 4\}, \{2, 4\}, \{3, 4\}, \{1, 5\}, \{1, 2, 5\}, \\ &\quad \{1, 3, 5\}, \{2, 3, 5\}, \{1, 4, 5\}, \{2, 4, 5\}\}. \end{aligned}$$

It is easily seen that the first is a cutset, while the second one is a collection with profile \mathbf{g} and not a cutset.

The next two propositions give us useful ways of determining whether a given vector $\mathbf{f} \in \Omega_n$ can be the profile of a cutset in $2^{[n]}$.

Proposition 2.4. *Let $\mathbf{f} \in \Omega_n$ and $0 \leq m_0 \leq n$. The canonical collection $\mathcal{C}(\mathbf{f}, m_0)$ defined above is a cutset if and only if $|\mathcal{E}_{m_0}^\uparrow| + |\mathcal{E}_{m_0}^\downarrow| \leq \binom{n}{m_0} + f_{m_0}$.*

Proof. If \mathcal{C} is a cutset, then the assertion follows as

$$|\mathcal{E}_{m_0}^\uparrow| + |\mathcal{E}_{m_0}^\downarrow| = |\mathcal{E}_{m_0}^\uparrow \cap \mathcal{E}_{m_0}^\downarrow| + |\mathcal{E}_{m_0}^\uparrow \cup \mathcal{E}_{m_0}^\downarrow| \leq |\mathcal{C}_{m_0}| + \binom{n}{m_0} \leq f_{m_0} + \binom{n}{m_0}.$$

On the other hand, if \mathcal{C} is not a cutset, then its profile must be exactly \mathbf{f} . Furthermore, as $\mathcal{E}_{m_0}^\downarrow$ is an initial segment and $\mathcal{E}_{m_0}^\uparrow$ is a last segment at level m_0 , their intersection has size greater than f_{m_0} and their union is all of $\binom{[n]}{m_0}$. Therefore, in this case we have $|\mathcal{E}_{m_0}^\uparrow| + |\mathcal{E}_{m_0}^\downarrow| = |\mathcal{E}_{m_0}^\uparrow \cap \mathcal{E}_{m_0}^\downarrow| + |\mathcal{E}_{m_0}^\uparrow \cup \mathcal{E}_{m_0}^\downarrow| > f_{m_0} + \binom{n}{m_0}$. \blacksquare

The next proposition shows the importance of the canonical collection. It further shows that the choice of m_0 is immaterial.

Proposition 2.5. *Let $\mathbf{f} \in \Omega_n$. The following are equivalent.*

1. \mathbf{f} is the profile of a cutset in $2^{[n]}$.
2. $\mathcal{C}(\mathbf{f}, m_0)$ is a cutset for some $0 \leq m_0 \leq n$.
3. $\mathcal{C}(\mathbf{f}, m_0)$ is a cutset for every $0 \leq m_0 \leq n$.

Proof. Clearly, if $\mathcal{C}(\mathbf{f}, m_0)$ is a cutset for some $0 \leq m_0 \leq n$, then there is a cutset in $2^{[n]}$ with profile \mathbf{f} . Therefore, it is enough to prove that if there is an $\mathcal{A} \subseteq 2^{[n]}$ which is a cutset with profile \mathbf{f} , then $\mathcal{C}(\mathbf{f}, m_0)$ is a cutset for every $0 \leq m_0 \leq n$.

For $m = 0, \dots, n$, let $\mathcal{A}_m = \{A \in \mathcal{A} \mid |A| = m\}$. We define \mathcal{B}_m^\uparrow and \mathcal{B}_m^\downarrow recursively by $\mathcal{B}_0^\uparrow = \{\emptyset\}$, $\mathcal{B}_m^\uparrow = \nabla(\mathcal{B}_{m-1}^\uparrow - \mathcal{A}_{m-1})$ for $m = 1, 2, \dots, n$, and $\mathcal{B}_n^\downarrow = \{[n]\}$, $\mathcal{B}_m^\downarrow = \Delta(\mathcal{B}_{m+1}^\downarrow - \mathcal{A}_{m+1})$ for $m = 0, 1, \dots, n-1$. Since \mathcal{A} is a cutset, we must have $\mathcal{B}_m^\downarrow \cap \mathcal{B}_m^\uparrow \subseteq \mathcal{A}_m$ for every $0 \leq m \leq n$.

Keeping the notations established prior to [Proposition 4](#), we now use downward induction on m to prove that $|\mathcal{B}_m^\downarrow| \geq |\mathcal{E}_m^\downarrow|$ and $|\mathcal{B}_m^\uparrow| \geq |\mathcal{E}_m^\uparrow|$ for every $0 \leq m \leq n$. Clearly $|\mathcal{B}_n^\downarrow| \geq |\mathcal{E}_n^\downarrow|$. Using, in order, the definition of \mathcal{B}_m^\downarrow , the Kruskal–Katona theorem, the triangle inequality, the inductive hypothesis, the fact that $\mathcal{C}_{m+1}^\downarrow \subseteq \mathcal{E}_{m+1}^\downarrow$, that $\mathcal{E}_{m+1}^\downarrow - \mathcal{C}_{m+1}^\downarrow$ is an initial segment at level $m+1$, and the definition of \mathcal{E}_m^\downarrow , we can write

$$\begin{aligned}
|\mathcal{B}_m^\downarrow| &= |\triangle (\mathcal{B}_{m+1}^\downarrow - \mathcal{A}_{m+1})| \\
&\geq |\triangle \mathcal{F}_{m+1}(|\mathcal{B}_{m+1}^\downarrow - \mathcal{A}_{m+1}|)| \\
&\geq |\triangle \mathcal{F}_{m+1}(|\mathcal{B}_{m+1}^\downarrow| - |\mathcal{A}_{m+1}|)| \\
&\geq |\triangle \mathcal{F}_{m+1}(|\mathcal{E}_{m+1}^\downarrow| - |\mathcal{C}_{m+1}^\downarrow|)| \\
&= |\triangle \mathcal{F}_{m+1}(|\mathcal{E}_{m+1}^\downarrow - \mathcal{C}_{m+1}^\downarrow|)| \\
&= |\triangle (\mathcal{E}_{m+1}^\downarrow - \mathcal{C}_{m+1}^\downarrow)| \\
&= |\mathcal{E}_m^\downarrow|,
\end{aligned}$$

as claimed. The assertion for $|\mathcal{B}_m^\uparrow|$ can be proved similarly.

Our assertion now follows from [Proposition 4](#), as

$$\begin{aligned}
|\mathcal{E}_{m_0}^\uparrow| + |\mathcal{E}_{m_0}^\downarrow| &\leq |\mathcal{B}_{m_0}^\uparrow| + |\mathcal{B}_{m_0}^\downarrow| = |\mathcal{B}_{m_0}^\uparrow \cap \mathcal{B}_{m_0}^\downarrow| + |\mathcal{B}_{m_0}^\uparrow \cup \mathcal{B}_{m_0}^\downarrow| \\
&\leq |\mathcal{A}_{m_0}| + |\mathcal{B}_{m_0}^\uparrow \cup \mathcal{B}_{m_0}^\downarrow| \leq f_{m_0} + \binom{n}{m_0}.
\end{aligned}$$

■

A quantitative version of [Proposition 5](#) can be formulated as follows.

Given positive integers K and m , there exist unique integers $a_m > a_{m-1} > \dots > a_t \geq t \geq 1$ such that

$$K = \binom{a_m}{m} + \binom{a_{m-1}}{m-1} + \dots + \binom{a_t}{t}.$$

This is called the *m-binomial representation* of K [[1](#), Theorem 7.2.1]. Using the *m*-binomial representation of K it is easy to describe the set numbered K in the squashed order on the *m*-th level of the Boolean lattice [[1](#), page 117].

For each positive integer m we define a (boundary) operator ∂_m [[11](#)] ([[6](#)] and [[9](#)] have other notations) on the integers as follows: If K is a positive integer with an *m*-binomial representation as above then

$$\partial_m(K) = \binom{a_m}{m-1} + \binom{a_{m-1}}{m-2} + \dots + \binom{a_t}{t-1},$$

and for non-positive K set $\partial_m(K) = 0$. Note that ∂_m is weakly increasing.

With this operator we can write

$$\triangle \mathcal{F}_m(K) = \mathcal{F}_{m-1}(\partial_m(K)), \quad \text{and} \quad \nabla \mathcal{L}_m(K) = \mathcal{L}_{m+1}(\partial_{n-m}(K)),$$

and the Kruskal–Katona theorem becomes $|\triangle \mathcal{B}| \geq \partial_m(|\mathcal{B}|)$ and $|\nabla \mathcal{B}| \geq \partial_{n-m}(|\mathcal{B}|)$ where \mathcal{B} is a collection of subsets of $[n]$ at level m .

Theorem 2.6. For a given $\mathbf{f} = (f_0, f_1, \dots, f_n) \in \Omega_n$, define

$$\mathbf{u}(\mathbf{f}) = (u_0, u_1, \dots, u_n) \quad \text{and} \quad \mathbf{v}(\mathbf{f}) = (v_0, v_1, \dots, v_n)$$

by:

$$\begin{aligned} u_0 &= 1 \text{ and } u_{m+1} = \partial_{n-m}(u_m - f_m) \text{ for } m=0, \dots, n-1, \\ v_n &= 1 \text{ and } v_{m-1} = \partial_m(v_m - f_m) \text{ for } m=1, \dots, n. \end{aligned}$$

Then the following are equivalent.

1. \mathbf{f} is the profile of a cutset in $2^{[n]}$.
2. $u_m + v_m - f_m \leq \binom{n}{m}$ for some $0 \leq m \leq n$.
3. $u_m + v_m - f_m \leq \binom{n}{m}$ for every $0 \leq m \leq n$.

Proof. (u_0, u_1, \dots, u_n) and (v_0, v_1, \dots, v_n) were chosen so that $|\mathcal{E}_m^\uparrow| = u_m$ and $|\mathcal{E}_m^\downarrow| = v_m$, so the statements follow from [Propositions 4 and 5](#). \blacksquare

Note that, as a special case when $f_0 = f_n = 0$, \mathbf{f} is the profile of a cutset in $2^{[n]}$ if and only if $v_0 = 0$ (or if and only if $u_n = 0$).

Remark. We can give a quantitative description of the canonical collection $\mathcal{C}(\mathbf{f}, m_0) = (\cup_{i=0}^{m_0} \mathcal{C}_i^\uparrow) \cup (\cup_{i=m_0+1}^n \mathcal{C}_i^\downarrow)$ as follows. If levels of $2^{[n]}$ are in the squashed order, then \mathcal{C}_i^\uparrow is the segment at level i starting with node numbered $\binom{n}{i} - u_i + 1$ and ending with $\min\{\binom{n}{i} - u_i + f_i, \binom{n}{i}\}$; \mathcal{C}_i^\downarrow is the segment at level i starting with node numbered $\max\{1, v_i - f_i + 1\}$ and ending with v_i .

3. Exact values: Proof of Theorem 1

We will use [Theorem 6](#) to determine the values of $g_n(m, l)$ for $l = m, m+1$, and $m+2$.

The case $l = m$ is obvious. For $l = m+1$, we can easily see that the canonical collection with profile $\mathbf{f} = (0, 0, \dots, 0, \binom{n-1}{m}, \binom{n-1}{m}, 0, \dots, 0)$ (where the nonzero components occur at levels m and $m+1$) consists of subsets of size m that do not contain n , together with subsets of size $m+1$ that contain n . This collection is a cutset. On the other hand, the canonical collection with profile $\mathbf{f}' = (0, 0, \dots, 0, \binom{n-1}{m}, \binom{n-1}{m} - 1, 0, \dots, 0)$ is not a cutset as there exists a maximal chain through the node $\{n-m, n-m+1, \dots, n\}$ which does not intersect the collection.

For the case $l = m+2$, the result is trivial for $m=0$ and hence assume that $m \geq 1$. Define $f = \sum_{j=0}^{m-1} \binom{n-2j-2}{m-j}$, and we need to prove that $g_n(m, l) = f+1$. We will show that $\mathbf{f} = (0, 0, \dots, 0, f+1, f, f, 0, \dots, 0)$ (where the nonzero

components occur at levels m , $m+1$, and $m+2$) is the profile of a cutset, but that $\mathbf{f}' = (0, 0, \dots, 0, f, f, f, 0, \dots, 0)$ is not the profile of a cutset.

Before starting the proof we need to establish a binomial identity which establishes a relationship between two *vertical* columns in the arithmetic (a.k.a. Pascal's) triangle. Let n be a positive integer, let $1 \leq m \leq n/2$, and let $0 \leq d \leq m-1$. Using downward induction on d (with base case $d=m-1$) it can be easily shown that

$$\sum_{j=0}^{m-d-1} \binom{n-2j-2}{m-j} + \sum_{j=0}^{m-d-1} \binom{n-2j-1}{m+2-j} + \binom{n-2m+2d}{d+2} = \binom{n}{m+2}.$$

In our computations below we rely on the case $d=0$ of this identity, namely:

$$f + \sum_{j=0}^{m-1} \binom{n-2j-1}{m+2-j} + \binom{n-2m}{2} = \binom{n}{m+2}.$$

We now use [Theorem 6](#) to show that \mathbf{f} is the profile of a cutset. Computing $\mathbf{v}(\mathbf{f})$ we get

$$\begin{aligned} v_i &= \binom{n}{i} && \text{for } i = n, n-1, \dots, m+2, \\ v_{m+1} &= \partial_{m+2} \left[\binom{n}{m+2} - f \right] \\ &= \partial_{m+2} \left[\sum_{j=0}^{m-1} \binom{n-2j-1}{m+2-j} + \binom{n-2m}{2} \right] \\ &= \sum_{j=0}^{m-1} \binom{n-2j-1}{m+1-j} + \binom{n-2m}{1}, \\ v_m &= \partial_{m+1} \left[\sum_{j=0}^{m-1} \binom{n-2j-1}{m+1-j} + \binom{n-2m}{1} - f \right] \\ &= \partial_{m+1} \left[\sum_{j=0}^{m-1} \binom{n-2j-2}{m+1-j} + \binom{n-2m}{1} \right] \\ &= \sum_{j=0}^{m-1} \binom{n-2j-2}{m-j} + \binom{n-2m}{0} \\ &= f + 1, \\ v_{m-1} &= \partial_m [f + 1 - f - 1] \\ &= 0, && \text{and} \end{aligned}$$

$$v_i = 0 \quad \text{for } i = m-1, m-2, \dots, 0,$$

hence \mathbf{f} is the profile of a cutset. As \mathbf{f}' differs from \mathbf{f} only in its m -th coordinate, we still have $v_m = f + 1$. Continuing, we get

$$\begin{aligned} v_{m-1} &= \partial_m[f + 1 - f] \\ &= \partial_m \left[\binom{m}{m} \right] \\ &= \binom{m}{m-1}, \quad \text{and} \\ v_i &= \binom{m}{i} \quad \text{for } i = m-1, m-2, \dots, 0. \end{aligned}$$

In particular, $v_0 = 1$, so \mathbf{f}' is not the profile of a cutset. ■

Remark. We can also construct a (non-canonical) cutset with f -vector $\mathbf{f} = (0, 0, \dots, 0, f+1, f, f, 0, \dots, 0)$ where $f = \sum_{j=0}^{m-1} \binom{n-2j-2}{m-j}$. We define $\mathcal{Q}_0 = \binom{[n-2]}{m}$ and for $j = 1, 2, \dots, m$ we let

$$\begin{aligned} \mathcal{Q}_j &= \left\{ A \cup \{n-1, n-3, \dots, n-2j+1\} \mid A \in \binom{[n-2j-2]}{m-j} \right\}, \\ \mathcal{R}_j &= \{Q \cup \{n-2j+2\} \mid Q \in \mathcal{Q}_{j-1}\}, \quad \text{and} \\ \mathcal{S}_j &= \{R \cup \{n-2j+1\} \mid R \in \mathcal{R}_{j-1}\}. \end{aligned}$$

Then it is rather straight forward to check that $\mathcal{Q} = \bigcup_{j=0}^m \mathcal{Q}_j$, $\mathcal{R} = \bigcup_{j=1}^m \mathcal{R}_j$, and $\mathcal{S} = \bigcup_{j=1}^m \mathcal{S}_j$ are collections of nodes at levels m , $m+1$, and $m+2$, respectively, $|\mathcal{Q}| = f+1$, $|\mathcal{R}| = f$, and $|\mathcal{S}| = f$, and that $\mathcal{Q} \cup \mathcal{R} \cup \mathcal{S}$ is a cutset, thus providing a cutset with f -vector $\mathbf{f} = (0, 0, \dots, 0, f+1, f, f, 0, \dots, 0)$. We omit the details but only briefly sketch the idea behind the construction: Note that \mathcal{Q}_0 consists of *all* subsets of size m that do not contain n or $n-1$. Likewise, \mathcal{R}_1 consists of all subsets of size $m+1$ that contain n but not $n-1$ and \mathcal{S}_1 consists of all subsets of size $m+2$ that contain n and $n-1$. Thus any maximal chain in $2^{[n]}$ that does not intersect $\mathcal{Q}_0 \cup \mathcal{R}_1 \cup \mathcal{S}_1$ will have subsets of size m , $m+1$ and $m+2$ that contain $n-1$ but not n . Now the poset of subsets of $[n]$ that contain $n-1$ but not n is isomorphic (as a poset) to $2^{[n-2]}$. We now restrict ourself to this poset and continue the construction recursively.

4. Bounds: Proof of Theorem 2

We need the following propositions.

Proposition 4.7. *For $m \leq l < n/2$ we have $g_n(m, n-m) \leq g_{n-1}(m, l)$.*

Proof. Let $\mathcal{A} \subseteq 2^{[n-1]}$ be a cutset with f -vector $(0, \dots, 0, g, \dots, g, 0, \dots, 0)$, where $g = g_{n-1}(m, l)$ and the nonzero entries are at levels $m \leq i \leq l$. Let $\overline{\mathcal{A}} = \{[n] \setminus A \mid A \in \mathcal{A}\}$. Note that by symmetry $\overline{\mathcal{A}}$ is a cutset in the poset \mathcal{Q} consisting of all subsets of $[n]$ that contain n . Now define $\mathcal{B} = \mathcal{A} \cup \overline{\mathcal{A}}$. Note that since $l < n-l$, \mathcal{B} has f -vector $(0, \dots, 0, g, \dots, g, 0, \dots, 0, g, \dots, g, 0, \dots, 0)$, where the nonzero entries are at levels $m \leq i \leq l$ and $n-l \leq i \leq n-m$. It now suffices to show that \mathcal{B} is a cutset in $2^{[n]}$.

Let $\mathcal{C} = C_0 \subset C_1 \subset \dots \subset C_n$ be a maximal chain in $2^{[n]}$ with $|C_i| = i$ for $0 \leq i \leq n$, and suppose, indirectly, that $\mathcal{B} \cap \mathcal{C} = \emptyset$. Since \mathcal{A} is a cutset in $2^{[n-1]}$ with support between levels m and l , we must have $n \in C_k$ for every $l \leq k \leq n$. Similarly, since $\overline{\mathcal{A}}$ is a cutset in \mathcal{Q} , we must have $n \notin C_k$ for every $0 \leq k \leq n-l$. This can only happen if $n-l < l$ which is a contradiction. ■

Proposition 4.8. *Suppose that $m+2 \leq l \leq n-m-1$ and $g_{n-1}(m+1, l-1) > \binom{n-2}{m}$. Then $g_n(m, l) > \binom{n-2}{m}$.*

Proof. Let $\mathbf{f} = (0, 0, \dots, 0, \binom{n-2}{m}, \binom{n-2}{m}, \dots, \binom{n-2}{m}, 0, \dots, 0)$, where the nonzero entries occur between levels m and l . Suppose, indirectly, that \mathbf{f} is the profile of a cutset in $2^{[n]}$. Then, by Proposition 5, the canonical collection $\mathcal{C}(\mathbf{f}, n)$ is such a cutset. If $C_i = \mathcal{C} \cap \binom{[n]}{i}$ are the level sets of the collection, then it is easy to see that $C_m = \binom{[n-2]}{m}$ and $C_{m+1} = \{A \cup \{n-1\} \mid A \in \binom{[n-2]}{m}\}$. Therefore $\cup_{i=m+2}^l C_i$ is a (canonical) cutset in $\{A \cup \{n\} \mid A \subseteq [n-1]\}$ of profile $(0, 0, \dots, 0, \binom{n-2}{m}, \binom{n-2}{m}, \dots, \binom{n-2}{m}, 0, \dots, 0)$, where the nonzero entries occur between levels $m+2$ and l . But $\{A \cup \{n\} \mid A \subseteq [n-1]\}$ is isomorphic to $2^{[n-1]}$, so we must have $g_{n-1}(m+1, l-1) \leq \binom{n-2}{m}$, a contradiction. ■

Proposition 4.9. *Define $\mathbf{f} = (f_0, f_1, \dots, f_n)$ as follows. Let $f_m = f_{m+1} = f_{n-m-1} = f_{n-m} = \binom{n-3}{m}$, $f_i = f_{n-i} = \binom{n+m-i-1}{m+1}$ for $i = m+2, m+3, \dots, \lfloor n/2 \rfloor$, and $f_i = 0$ otherwise. Then \mathbf{f} is not the profile of a cutset in $2^{[n]}$.*

Proof. For $\mathbf{v}(\mathbf{f})$ we get the following.

$$v_i = \binom{n}{i} \quad \text{for } i = n, n-1, \dots, n-m,$$

$$\begin{aligned}
v_{n-m-1} &= \partial_{n-m} \left[\binom{n}{n-m} - \binom{n-3}{n-m-3} \right] \\
&= \partial_{n-m} \left[\binom{n-1}{n-m} + \binom{n-2}{n-m-1} + \binom{n-3}{n-m-2} \right] \\
&= \binom{n-1}{n-m-1} + \binom{n-2}{n-m-2} + \binom{n-3}{n-m-3}, \\
v_{n-m-2} &= \partial_{n-m-1} \left[\binom{n-1}{n-m-1} + \binom{n-2}{n-m-2} \right. \\
&\quad \left. + \binom{n-3}{n-m-3} - \binom{n-3}{n-m-3} \right] \\
&= \partial_{n-m-1} \left[\binom{n-1}{n-m-1} + \binom{n-2}{n-m-2} \right] \\
&= \binom{n-1}{n-m-2} + \binom{n-2}{n-m-3}, \\
v_{n-m-3} &= \partial_{n-m-2} \left[\binom{n-1}{n-m-2} + \binom{n-2}{n-m-3} - \binom{n-3}{n-m-4} \right] \\
&= \partial_{n-m-2} \left[\binom{n-1}{n-m-2} + \binom{n-3}{n-m-3} \right] \\
&= \binom{n-1}{n-m-3} + \binom{n-3}{n-m-4}, \\
v_{n-m-4} &= \partial_{n-m-3} \left[\binom{n-1}{n-m-3} + \binom{n-3}{n-m-4} - \binom{n-4}{n-m-5} \right] \\
&= \partial_{n-m-3} \left[\binom{n-1}{n-m-3} + \binom{n-4}{n-m-4} \right] \\
&= \binom{n-1}{n-m-4} + \binom{n-4}{n-m-5},
\end{aligned}$$

and continuing as above, we get $v_k = \binom{n-1}{k} + \binom{k+m}{k-1}$ for every $k \geq \lfloor (n-1)/2 \rfloor$.

Now note that since \mathbf{f} is symmetrical, for $\mathbf{v}(\mathbf{f})$ and $\mathbf{u}(\mathbf{f})$ we have $v_i = u_{n-i}$ for every $0 \leq i \leq n$. Therefore

$$\begin{aligned}
v_{\lfloor n/2 \rfloor} + u_{\lfloor n/2 \rfloor} &= v_{\lfloor n/2 \rfloor} + v_{\lceil n/2 \rceil} \\
&= \binom{n-1}{\lfloor n/2 \rfloor} + \binom{\lfloor n/2 \rfloor + m}{\lfloor n/2 \rfloor - 1} + \binom{n-1}{\lceil n/2 \rceil} + \binom{\lceil n/2 \rceil + m}{\lceil n/2 \rceil - 1}
\end{aligned}$$

$$= \binom{n}{\lfloor n/2 \rfloor} + f_{\lfloor n/2 \rfloor} + f_{\lceil n/2 \rceil},$$

thus \mathbf{f} cannot be the profile of a cutset by [Theorem 6](#). ■

Proof of [Theorem 2](#) and [Corollary 3](#). The upper bound in 1 follows from [Theorem 1](#) part 3 since $g_n(m, l)$ is weakly decreasing with l . The upper bound in 2 is a consequence of [Proposition 7](#) (take $l = m + 2$ assuming $n > 2m + 4$) and [Theorem 1](#) part 3. [Proposition 9](#) implies that $g_n(m, n - m) > \min\{\binom{n-3}{m}, \binom{\lfloor n/2 \rfloor + m - 1}{m+1}\}$. It is an elementary exercise to check that this minimum is $\binom{n-3}{m}$ for $n \gg m$ (e.g. if $n \geq 3^{m+1}(m+1)$), establishing the lower bound in 2. Finally, the lower bound in 1 is a consequence of [Proposition 8](#) taking $l = n - m - 1$, and noting that $\binom{n-4}{m+1} \geq \binom{n-2}{m}$ for $n \gg m$ (e.g. if $n \geq 8(m+1)$) and the lower bound in 2.

The cases of $l = 1, 2$ (and 3) of [Corollary 3](#) follow from [Theorem 1](#). The assertions for $3 \leq l \leq n - 1$ follow from [Theorem 2](#) if $n \geq 18$ (see above). The cases $5 \leq n \leq 17$ were checked directly using [Theorem 6](#) (and a simple computer program). ■

We close this section by proving a partial complement to [Proposition 7](#).

Proposition 4.10. *For $n/2 - 1 < l \leq n - m - 2$ we have $g_n(m, n - m) \geq g_{n-1}(m, l)$.*

Proof. Consider the canonical collection $\mathcal{C} = \mathcal{C}(\mathbf{f}, l)$ where $\mathbf{f} = (f_0, f_1, \dots, f_n)$ is defined by $f_i = g_n(m, n - m)$ for $m \leq i \leq n - m$ and 0 otherwise. Then \mathcal{C} is a cutset in $2^{[n]}$. Write $\mathcal{D}_1 = \{C \in \mathcal{C} \mid |C| \leq l\}$ and $\mathcal{D}_2 = \{C \in \mathcal{C} \mid |C| \geq l + 1\}$.

Suppose, indirectly, that $g_n(m, n - m) < g_{n-1}(m, l)$. Then \mathcal{D}_1 is not a cutset in $2^{[n-1]}$, in particular, $\mathcal{D}_1 \subseteq 2^{[n-1]}$. Since $l \geq n - 1 - l$, by symmetry \mathcal{D}_2 is disjoint from $2^{[n-1]}$. Therefore $\mathcal{D}_1 \cup \mathcal{D}_2 = \mathcal{C}$ is not a cutset in $2^{[n]}$, a contradiction. ■

5. Examples and conjectures

The following table has the exact values of $g_n(m, l)$ for $n = 100$, $m = 4$ and every $4 \leq l \leq 96$. These were found using [Theorem 6](#). We have also given the 4-binomial representation of $g_n(m, l)$ in the third column.

l	$g_{100}(4, l)$	4-binomial representation of $g_{100}(4, l)$
4	3,921,225	$\binom{100}{4}$
5	3,764,376	$\binom{99}{4}$
6	3,759,624	$\binom{98}{4} + \binom{96}{3} + \binom{94}{2} + \binom{93}{1}$
7	3,759,526	$\binom{98}{4} + \binom{96}{3} + \binom{93}{2} + \binom{88}{1}$
$8 \leq l \leq 95$	3,759,525	$\binom{98}{4} + \binom{96}{3} + \binom{93}{2} + \binom{87}{1} = \binom{100}{4} - \binom{100}{3}$
96	3,607,527	$\binom{97}{4} + \binom{95}{3} + \binom{92}{2} + \binom{86}{1} = \binom{99}{4} - \binom{99}{3}$

Theorem 1 gives the exact value for $4 \leq l \leq 6$. For $6 \leq l \leq 95$, **Theorem 2** gives $3,612,280 = \binom{98}{4} < g_{100}(4, l) \leq \binom{98}{4} + \binom{96}{3} + \binom{94}{2} + \binom{92}{1} + \binom{90}{0} = 3,759,624$. Finally, for $l = 96$, **Theorem 2** gives $3,464,840 = \binom{97}{4} < g_{100}(4, 96) \leq \binom{97}{4} + \binom{95}{3} + \binom{91}{2} + \binom{89}{1} + \binom{80}{0} = 3,607,625$.

From these and other similar tables we see how $g_n(m, l)$ decreases as l increases from m to $n - m$. Namely, we observe that the decrease is largest from level m to level $m + 1$ and from level $n - m - 1$ to level $n - m$, quite modest between level $m + 2$ and level $2m$, and that, rather strikingly, $g_n(m, l)$ is constant between levels $l = 2m$ and $l = n - m - 1$. In fact, we have the following conjectures.

Conjecture 5.11. For $n \gg m$ we have

1. $g_n(m, l) = \binom{n}{m} - \binom{n}{m-1}$ for every $l = 2m, 2m + 1, \dots, n - m - 1$, and
2. $g_n(m, n - m) = \binom{n-1}{m} - \binom{n-1}{m-1}$.

Note that the m -binomial representation of $\binom{n}{m} - \binom{n}{m-1}$ starts with $\binom{n-2}{m} + \binom{n-4}{m-1}$ (and $\binom{n-1}{m} - \binom{n-1}{m-1}$ starts with $\binom{n-3}{m} + \binom{n-5}{m-1}$) when $n \gg m$ (cf. **Theorem 2**). **Corollary 3** proves our conjectures for $m = 1$, and **Theorem 2** establishes that **Conjecture 11** provides an upper bound for $m = 2$ since $\binom{n}{2} - \binom{n}{1} = \binom{n-2}{2} + \binom{n-3}{1}$ (and $\binom{n-1}{2} - \binom{n-1}{1} = \binom{n-3}{2} + \binom{n-4}{1}$).

According to **Conjecture 11**, we have $g_n(m, n - m) = g_{n-1}(m, n - m - 2) = g_{n-1}(m, l)$ for $n \gg m$ and $2m \leq l \leq n - m - 2$. **Proposition 7** establishes $g_n(m, n - m) \leq g_{n-1}(m, l)$ for $l < n/2$, while **Proposition 10** proves the other direction for $n/2 - 1 < l \leq n - m - 2$ (yielding equality when $l = (n - 1)/2$).

Note. In a subsequent paper [2], the above conjectures have been somewhat refined and related to other conjectures about the width of cutsets in the truncated Boolean lattice.

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